Determining the Intrinsic Structure of Public Software Development History

Antoine Pietri
antoine.pietri@inria.fr
Inria
Paris, France

Guillaume Rousseau
guillaume.rousseau@univ-paris-diderot.fr
Université de Paris and Inria
Paris, France

Stefano Zacchiroli
zack@irif.fr
Université de Paris and Inria
Paris, France

ABSTRACT

Background. Collaborative software development has produced a wealth of version control system (VCS) data that can now be analyzed in full. Little is known about the intrinsic structure of the entire corpus of publicly available VCS as an interconnected graph. Understanding its structure is needed to determine the best approach to analyze it in full and to avoid methodological pitfalls when doing so.

Objective. We intend to determine the most salient network topology properties of public software development history as captured by VCS. We will explore: degree distributions, determining whether they are scale-free or not; distribution of connect component sizes; distribution of shortest path lengths.

Method. We will use Software Heritage—which is the largest corpus of public VCS data—compress it using webgraph compression techniques, and analyze it in-memory using classic graph algorithms. Analyses will be performed both on the full graph and on relevant subgraphs.

Limitations. The study is exploratory in nature; as such no hypotheses on the findings is stated at this time. Chosen graph algorithms are expected to scale to the corpus size, but it will need to be confirmed experimentally. External validity will depend on how representative Software Heritage is of the software commons.

KEYWORDS

source code, version control system, network topology, graph structure, statistical mechanics

ACM Reference Format:

1 INTRODUCTION

The rise in popularity of Free/Open Source Software (FOSS) and collaborative development platforms [10] over the past decades has made publicly available a wealth of software source code artifacts (source code files, commits with all associated metadata, tagged released, etc.), which have in turn benefited empirical software engineering (ESE) and mining software repository (MSR) research. Version control systems (VCS) in particular have been frequently analyzed [9] due to the rich view they provide on software evolution and their ease of exploitation since the advent of distributed version control systems.

Only very recently systematic initiatives [1, 6, 12] have been established to gather as much public VCS data as possible in a single logical place, enabling ESE research at the scale of, ideally, all publicly available source code artifacts—i.e., our entire software commons [11]. In this paper we describe how we will conduct the first systematic exploratory study on the intrinsic structure of source code artifacts found in version control systems.

As corpus we will use Software Heritage and its dataset [1, 15], which is the largest and most diverse collection of source code artifacts, spanning more than 5 billion unique source code files and 1 billion unique commits, collected from more than 80 million software projects.
The data model of our corpus is shown in Figure 1. It is a Merkle DAG [14], fully deduplicated, linking together in a single giant graph the development histories of all publicly available VCSs. Actual source code files are represented as blob nodes, their grouping in source code trees as directory nodes; commit nodes are linked together by ancestry, supporting forks and N-ary merges; releases (or “tags”) are also supported; the full states of VCSs is recorded periodically (similarly to what the Internet Archive wayback machine does for web pages) and associated to where it has been observed (a software origin, identified by an URL). We refer the reader to [15] for full details on the data model.

The study we will conduct will assess the most salient network topology properties [2] of the Software Heritage corpus as a graph. Such a systematic analysis is still missing in the literature and is needed for further empirical software engineering research for several reasons:

a) **Determine the most appropriate large-scale analysis approach.** Most “large-scale” studies of VCSs fall short of the full body of publicly available source code artifacts and either resort to random sampling or focus on popular repositories. This is a potential source of bias, but is understandable for practical reasons. To enable studies on the entire software commons we need, in addition to platforms [1, 12], an understanding of its intrinsic structure, to choose the most appropriate large-scale analysis approach [7, 16] depending on the study needs. For instance, if the graph is easy to partition into loosely connected components, then a scale-out approach with several compute nodes holding in memory graph quasi-partitions would be best; if the graph is highly connected a scale-up approach relying on graph compression [3] would be preferable. Similarly, knowing that most nodes are part of a single gigantic connect component (CC) would help in avoiding algorithmic approaches with high complexity on the size of the largest CC.

b) **Avoid methodological pitfalls.** The same understanding is needed to avoid making strong assumptions on what constitutes “typical” VCS data. These pitfalls have been warned against since the early days of GitHub [10], but that have not been quantified yet. The extent to which repositories on popular forges correspond to “well behaved” development repositories, as opposed to be outliers that are not used for software development or are built just to test the limits of hosting platforms or VCS technology is unknown. In our experience GitHub alone contains repositories with very weird artifacts (commits with one million parents or mimicking bitcoin mining in their IDs, the longest possible paths, bogus timestamps [16], etc.). How many statistically relevant outliers of this kind exist is unknown and needs to be documented as reference knowledge to help researchers in the interpretation of their empirical findings.

c) **Improve our understanding of our daily objects of study.** More generally, in empirical software engineering we are collectively studying artifacts that naturally emerge from the human activity of software development. As it is commonplace in other sciences (and most notably physics), we want to study the intrinsic network properties of the development history of our software commons just because the corpus exists, it is available, and it is challenging to do so. The resulting findings might be also practically useful, but in spite of that we will obtain a more deep understanding of the nature of objects we study daily than what is known today.

### Table 1: Corpus size as a graph (release: 2018-09-25).

<table>
<thead>
<tr>
<th>Nodes</th>
<th>Edges</th>
</tr>
</thead>
<tbody>
<tr>
<td>origins (ori)</td>
<td>ori→snp</td>
</tr>
<tr>
<td>snapshots (snp)</td>
<td>snp→cmt</td>
</tr>
<tr>
<td>releases (rel)</td>
<td>cmt→cmt</td>
</tr>
<tr>
<td>commits (cmt)</td>
<td>cmt→dir</td>
</tr>
<tr>
<td>directories (dir)</td>
<td>dir→dir</td>
</tr>
<tr>
<td>files</td>
<td>dir→file</td>
</tr>
<tr>
<td></td>
<td>≈ 11 B</td>
</tr>
<tr>
<td></td>
<td>≈ 165 B</td>
</tr>
</tbody>
</table>

2 **RESEARCH QUESTIONS**

Specifically we will perform an exploratory study, with no predetermined hypotheses, and answer the following research questions:

RQ1 What is the distribution of indegrees, outdegrees and local clustering of the public VCS history graph? Which laws do they fit? How do such distributions vary across the different graph layers—file system layer (files + directories) v. history layer (commits + releases) v. origin layer.

RQ2 What is the distribution of connected component sizes for the public VCS history graph? How does it vary across graph layers?

RQ3 What is the distribution of shortest path lengths from roots to leaves in the recursive layers (commits and directories) of public VCS history graph.

3 **VARIABLES**

We define the following subsets of the starting corpus:

- **full corpus**: the entire Software Heritage graph dataset
- **filesystem layer**: full corpus subset consisting of file and directory nodes only, and edges between them
- **history layer**: full corpus subset consisting of commit and releases only, and edges between them
- **commit layer**: subset of the history layer consisting of commit nodes only, and edges between them
- **hosting layer**: full corpus subset consisting of origins and snapshot nodes only, and edges between them

For each corpus we will measure the following variables:

- **indegree distribution**: for each node the number of edges pointing to it
- **outdegree distribution**: for each node the number of edges starting from it
- **local undirected clustering distribution**: for each node the number of edges between nodes pointing to or from it (= local clustering coefficient without dividing it by the number of possible egdes between its neighbors)
- **CC size distribution**: the size, in number of nodes, of each connected component in the underlying undirected graph of the input corpus
- **(for the filesystem and commit layers only) path length distribution**: distribution of the length of shortest paths between all root (indegree=0) nodes and all leaf (outdegree=0) nodes
4 MATERIAL

Dataset. We intend to use the Software Heritage graph dataset [15] as corpus for the planned experiments. The main reasons for this choice are that: (1) it is the largest dataset about publicly software development history (see Table 1 for its size characteristics as a graph); (2) it is available as open data in various formats, including a simple nodes/edges graph representation; (3) it has been chosen as topic for the MSR 2020 Mining Challenge, hence we expect that by the time experiments will be run the body of related work around it will be substantial.

Software. We will use swb-graph [3] and WebGraph [4] to produce and exploit compressed graph representations of the input corpus (see Section 5 for details). Statistical analyses will be performed using popular SciPy [13] components (NumPy, Pandas, . . . ). All developed custom code and derived data will be released as a complete replication package.

Hardware. We expect to be able to run all experiments on a single server equipped with 24 Intel Xeon 2.20 GHz CPUs and 750 GB of RAM, which is already available to us.

5 EXECUTION PLAN

We will follow the experiment protocol described below.

a. Retrieve the most recent version of the Software Heritage graph dataset [15] available, falling back to the one available at the time of writing (2018-09-25) if no newer releases have been published.

b. Compress the textual node/graph representation of the full dataset to a compact representation of its structure (which ignores all node metadata except node types) using swb-graph. Repeat for each subgraph of interest (see Section 3 for details). This step is expected to take about a week for the full graph and proportionally less time for the various subgraphs.

c. (RQ1) Compute indegrees, outdegrees and local clustering for all nodes in the graph, for all relevant (sub)graphs, exploiting the WebGraph [4] API with custom Java code.

d. (RQ2) Compute all connected components on all relevant (sub)graphs using well known algorithms [8] that can be implemented with custom Java code as visits (BFS or DFS) on the WebGraph graph representations. A full graph visit is expected to take a few hours, scaling down linearly for subgraphs. [3]

e. (RQ3) For each subgraph of interest (e.g., the file system layer composed only of files and directory nodes), for all root nodes (e.g., source code root directories), create a shortest path spanning tree to all leaves (e.g. file or directory nodes with no children). Then export all path lengths. This can be implemented with custom Java code realizing Dijkstra’s algorithm on top of the WebGraph API. We have no precise estimate of how long this step will take: it will take more than a single full graph visit, due to paths that will be re-explored, but it should remain manageable. Thread-based parallelization is an option in case we need to speed up this step.

6 ANALYSIS PLAN

6.1 Descriptive statistics

We will analyze the raw data obtained from the execution plan according to the following analysis protocol.

a. (RQ1) We will plot indegree/outdegree/local clustering distributions as histograms, one histogram per distribution and per (sub)graph (see Section 3 for details).

   • One figure per distribution and graph will be produced, showing both the raw histogram (i.e., for each degree/clustering value observed, the number of nodes with that value) as well as the cumulative distribution function (CDF) given by the number of nodes whose degree/clustering is greater than or equal to a given value. One or more scales among lin-lin, lin-log, log-lin, log-log scales will be used, depending on presentation needs.
   • Each distribution will be qualitatively discussed by comparing it with a generic distribution functions (power law, exponential law, Poisson law).
   • In case they arise, the presence of outliers or other statistical anomalies with respect to more regular regions of the distribution will be highlighted (e.g., change of law, change of slope, . . . ), and may lead to specific investigations to determine their nature.

b. (RQ1) The nature of the tail of each distribution will be systematically analyzed and discussed.

   • The first observed criterion will be the amplitude of the range of degree values, expressed in decades.
   • Then, we will use the discrete maximum likelihood estimator (MLE) to determine the scaling parameter [5, Eq. 3.7].

\[
\alpha(d_{\text{min}}) = 1 + n \sum_{i=1}^{n} \ln \frac{d_i}{d_{\text{min}}}^{-1}
\]

where \(d_i, i = 1, \ldots, n\) are the observed degree values such that \(d_i \geq d_{\text{min}}\).

Without prejudging, nor speculating whether the distributions will match power laws or not, we limit ourselves to the first step of the methodology proposed in [5, Box 1], using the above estimator which depends on an arbitrary degree threshold \(d_{\text{min}}\), beyond which the behaviour of the distribution is ignored, like a probe through the whole range of degrees.

The plot is displayed in log-lin scale over the degree range. Lacking precise information on the results at this stage, as well as knowledge about the impact of possible atypical events in the observed distributions, it is premature to plan to implement the subsequent steps of the aforementioned methodology.

c. (RQ2) We will display connected component size, focusing on VCS origins that contain at least one commit, as a set of histograms.

   • As above, both raw distribution functions and cumulative distributions will be displayed.
   • To characterize the aggregation process into connected components as nodes of different types, or at different depths, are added to the corpus:

\[\bar{\alpha}(d_{\text{min}}) = 1 + n \sum_{i=1}^{n} \ln \frac{d_i}{d_{\text{min}}}^{-1}\]
6.2 Practical significance of the findings

RQ1’s answers will provide information on how the public VCS graph compares with other naturally occurring graphs in collaboration, such as the social network graphs or the graph of the Web. While seemingly only of theoretical significance, degree distributions and derived properties such as graph density directly impact on the practical exploitablety of graphs of this scale. For instance graphs with fat-tailed degree distributions compress better than others. We know from previous work [3, 16] how well the full VCS graph compress, but a systematic study of related properties for the subgraphs we intend to address here is still missing.

Findings related to RQ2 will directly tell how to best approach full-scale analyses of the various corpuses. If, for instance, it will turn out that 80% of the nodes are part of a gigantic component, than distributed approaches will be generally difficult to implement, due to entanglement and the pseudo-random nature of identifiers in Merkle DAG. Opposite considerations will apply for a more uniform distribution of CC sizes.

Finally, RQ3’s answers will tell what’s the average path depths in the recursive layers of the public VCS corpus (commit and filesystem layers), as a pair of histograms.

- As before, both raw and cumulative distribution functions will be displayed.

7 LIMITATIONS

Exploratory nature of the study. The study we propose is exploratory in nature, hence we state no hypotheses on the findings at this stage. This is intended, as the intrinsic structure of the public VCS corpus has never been characterized before at the extent we propose.

Algorithmic feasibility. With its 10^12 nodes and 160^12 edges the corpus we plan to study is a substantial graph for graph practitioners standards; not as big as the graph of the Web, but significantly larger than most benchmarks used in the field. At this scale, algorithms with super-linear complexity are generally considered non practically applicable.

Our algorithmic approaches for RQ1 and RQ2 have linear complexity; the approach for RQ3 is super linear, but we are restricting it to selected subgraphs and we expect it to exhibit enough sharing/caching to be practically treatable. We do not expect it to be the case but if, due to either algorithmic or technological considerations, any specific sub-experiment will turn out to be not practically feasible, we will resort to uniform random sampling.

External validity. Due to organic crawling lag, Software Heritage does not capture the full extent of publicly available VCSs. Hence we do not claim being able to characterize the intrinsic structure of the entire history of publicly available software development. The chosen dataset is nonetheless the best publicly accessible option available today to researchers. It is also well representative of the most popular development forge(s) in use today. As such we expect that the network topology findings of this study will provide useful insights to researchers and practitioners in the field.

REFERENCES